

# Pushforwards in Inverse Homotopical Diagrams



Krzysztof Kapulkin and [Yufeng Li](#)



## Background

Krzysztof **Kapulkin**, and Peter LeFanu **Lumsdaine**. **The homotopy theory of type theories**. *Advances in Mathematics*, 2018.



## Morita equivalence

The **category of models of ITT** admits a **cofibrantly-generated model structure**



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## Path object construction

- ▶ The path object model for each  $\mathbb{C} \in \text{Mod}_{\text{ITT}}$  is an inverse homotopical diagram category with Reedy fibrations



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## Path object construction

- ▶ The path object model for each  $\mathbb{C} \in \text{Mod}_{\text{ITT}}$  is an inverse homotopical diagram category with Reedy fibrations
- ▶ Must ensure the path object model supports various logical structures: most challenging of which is the  $\Pi$ -types (i.e. pushforwards)



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Krzysztof **Kapulkin**, and Peter LeFanu **Lumsdaine**. **Homotopical inverse diagrams in categories with attributes**. *Journal of Pure and Applied Algebra*, 2021.

Marcelo **Fiore**, Krzysztof **Kapulkin**, and Y.L. **Logical structures in inverse functor categories**. *arxiv:2410.11728*, 2024.



## Main result

Krzysztof **Kapulkin** and Y.L. **Pushforwards in Inverse Homotopical Diagrams**. *arxiv:2506.04472*, 2025.





## Definition

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If  $\mathbb{C}$  is a model of logic then one would like  $\mathbb{C}^{\mathcal{I}}$  to be a model of logic, for inverse categories  $\mathcal{I}$ .



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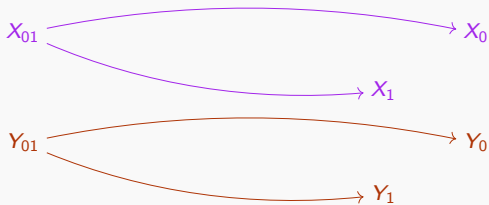
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- ▶  $\mathbb{C}^{\mathcal{I}}$  can often **validate or refute additional axioms**:  $\text{Set}^{\rightarrow}$  refutes **excluded middle**
- ▶ The **path object** for each  $\mathbb{C} \in \text{Mod}_{\text{ITT}}$  is an **inverse homotopical diagram category** with **Reedy fibrations** in the **Kapulkin–Lumsdaine** model structure



## Constructing spans inductively

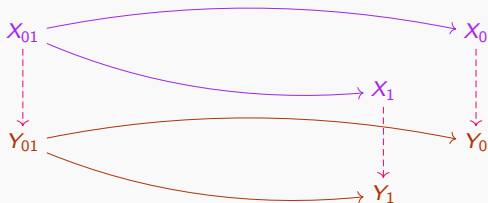
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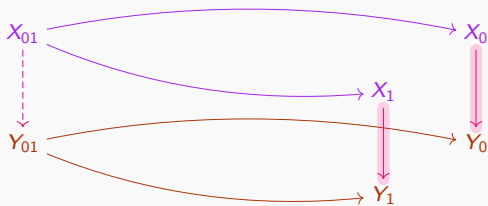
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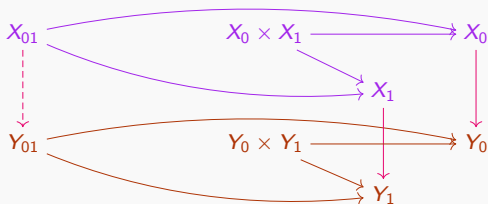
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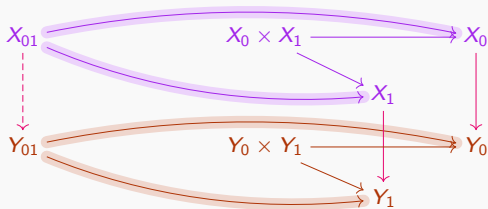
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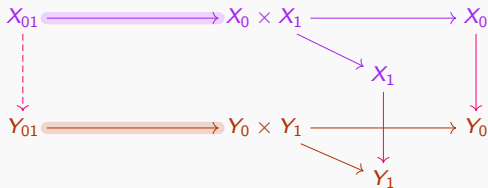
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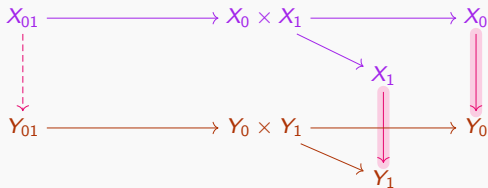
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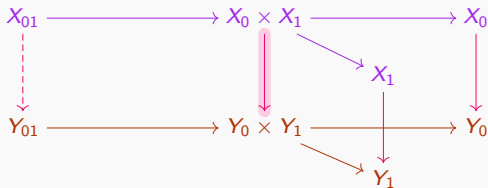
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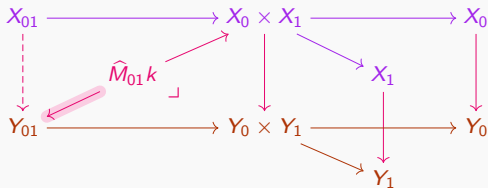
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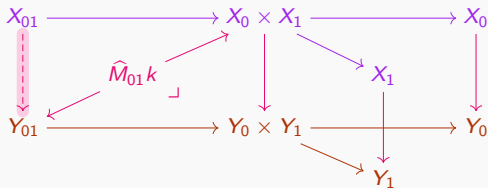
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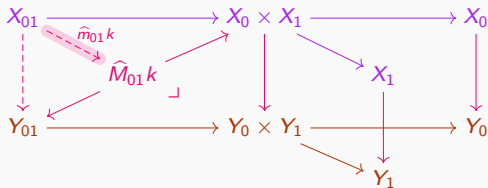
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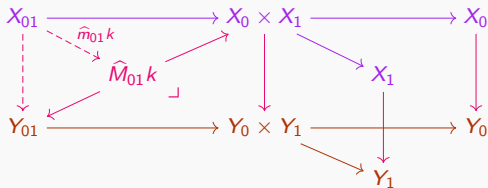
Suppose we want to **construct**  $k: X \rightarrow Y$  in  $\mathbb{C}^{0 \leftarrow 01 \rightarrow 1}$  from **components at 0, 1**.  
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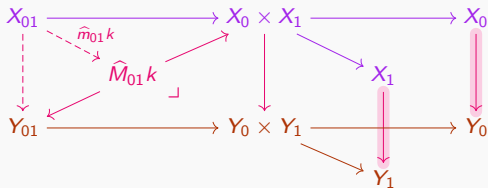






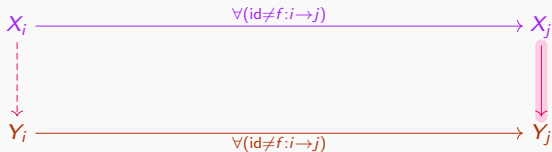
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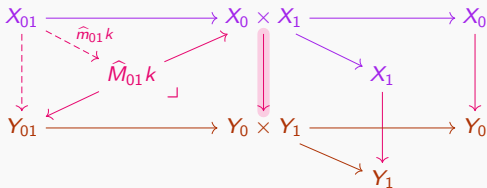
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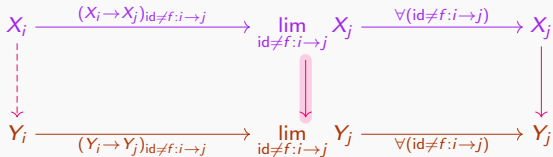
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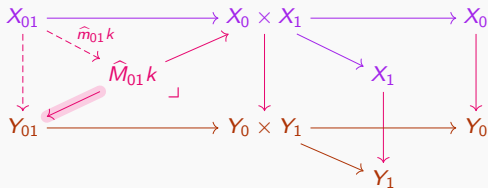
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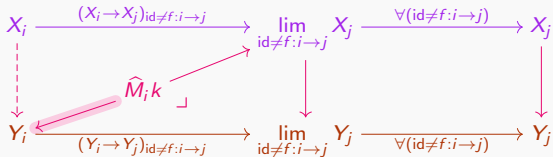
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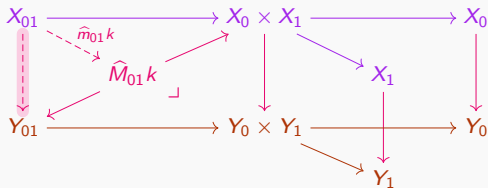
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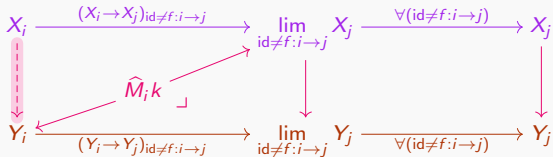
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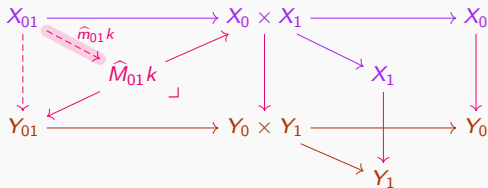
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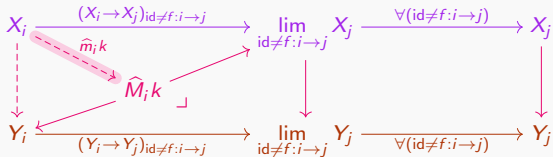
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## Constructing diagrams inductively

Generally, for **inverse**  $\mathcal{I}$ , constructing the  $i$ -th **component** of **construct**  $k: X \rightarrow Y$  from those at **lower degrees**  $j$  amounts to a map  $\widehat{m}_i k: X_i \rightarrow \widehat{M}_i k$





## Definition

Suppose  $\mathbb{C}$  is equipped with a class of **pullback-stable fibrations**.

$$\begin{array}{ccccc}
 X_i & \xrightarrow{(X_i \rightarrow X_j)_{id \neq f: i \rightarrow j}} & \lim_{id \neq f: i \rightarrow j} X_j & \xrightarrow{\forall (id \neq f: i \rightarrow j)} & X_j \\
 \downarrow \widehat{m}_{i,k} & \nearrow & \downarrow & & \downarrow \\
 Y_i & \xrightarrow{(Y_i \rightarrow Y_j)_{id \neq f: i \rightarrow j}} & \lim_{id \neq f: i \rightarrow j} Y_j & \xrightarrow{\forall (id \neq f: i \rightarrow j)} & Y_j \\
 & & \lrcorner & & 
 \end{array}$$

The diagram illustrates a commutative square with a pullback. The top row shows a map from  $X_i$  to the limit  $\lim_{id \neq f: i \rightarrow j} X_j$  labeled  $(X_i \rightarrow X_j)_{id \neq f: i \rightarrow j}$ , followed by a map to  $X_j$  labeled  $\forall (id \neq f: i \rightarrow j)$ . The bottom row shows a map from  $Y_i$  to the limit  $\lim_{id \neq f: i \rightarrow j} Y_j$  labeled  $(Y_i \rightarrow Y_j)_{id \neq f: i \rightarrow j}$ , followed by a map to  $Y_j$  labeled  $\forall (id \neq f: i \rightarrow j)$ . A vertical dashed arrow from  $X_i$  to  $Y_i$  is labeled  $\widehat{m}_{i,k}$ . A diagonal arrow from  $X_i$  to the limit  $\lim_{id \neq f: i \rightarrow j} Y_j$  is labeled  $\widehat{M}_{i,k}$ . A right-angle symbol  $\lrcorner$  is placed at the bottom right of the square formed by  $\widehat{M}_{i,k}$  and the bottom row.

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Suppose  $\mathcal{C}$  is equipped with a class of **pullback-stable fibrations**.

For inverse  $\mathcal{I}$ , a map  $k: X \rightarrow Y$  in  $\mathcal{C}^{\mathcal{I}}$  is a **Reedy fibration** when each  $\widehat{m}_i k: X_i \rightarrow \widehat{M}_i k$  is a **fibration** in  $\mathcal{C}$ .

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The diagram illustrates the definition of a Reedy fibration. It shows a commutative square of maps between objects in a category  $\mathcal{C}$ . The top row consists of  $X_i$ ,  $\lim_{\text{id} \neq f: i \rightarrow j} X_j$ , and  $X_j$ . The bottom row consists of  $Y_i$ ,  $\lim_{\text{id} \neq f: i \rightarrow j} Y_j$ , and  $Y_j$ . The left vertical arrow is  $\widehat{m}_i k$ , the right vertical arrow is the identity, and the bottom horizontal arrow is the identity. The top horizontal arrow is labeled  $(X_i \rightarrow X_j)_{\text{id} \neq f: i \rightarrow j}$ , and the bottom horizontal arrow is labeled  $(Y_i \rightarrow Y_j)_{\text{id} \neq f: i \rightarrow j}$ . The diagonal arrow from  $X_i$  to  $\lim_{\text{id} \neq f: i \rightarrow j} Y_j$  is labeled  $\widehat{M}_i k$ . A right-angle symbol is shown at the vertex  $\lim_{\text{id} \neq f: i \rightarrow j} Y_j$ .



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The diagram illustrates the commutativity of the square formed by the maps  $X_i \rightarrow \lim_{id \neq f: i \rightarrow j} X_j$ ,  $\lim_{id \neq f: i \rightarrow j} X_j \rightarrow X_j$ ,  $X_i \rightarrow Y_i$ , and  $Y_i \rightarrow \lim_{id \neq f: i \rightarrow j} Y_j$ . The map  $\widehat{m}_i k$  is shown as a dashed arrow from  $X_i$  to  $\widehat{M}_i k$ , and the map  $\widehat{M}_i k$  is shown as a solid arrow from  $\widehat{M}_i k$  to  $\lim_{id \neq f: i \rightarrow j} X_j$ . A right-angle symbol is placed at the vertex  $\widehat{M}_i k$  to indicate that the square is a pullback.

## Definition

A category with weak equivalences is a category  $\mathbb{C}$  equipped with a **replete wide subcategory**  $\mathcal{W}_{\mathbb{C}} \hookrightarrow \mathbb{C}$  closed under **two-out-of-three** called the weak equivalences.

A homotopical diagram  $(\mathcal{I}, \mathcal{W}_{\mathcal{I}}) \rightarrow (\mathbb{C}, \mathcal{W}_{\mathbb{C}})$  is a diagram  $\mathcal{I} \rightarrow \mathbb{C}$  preserving weak equivalences.

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For inverse  $\mathcal{I}$ , a map  $k: X \rightarrow Y$  in  $\mathbb{C}^{\mathcal{I}}$  is a **Reedy fibration** when each  $\widehat{m}_i k: X_i \rightarrow \widehat{M}_i k$  is a **fibration** in  $\mathbb{C}$ .

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 X_i & \xrightarrow{(X_i \rightarrow X_j)_{id \neq f: i \rightarrow j}} & \lim_{id \neq f: i \rightarrow j} X_j & \xrightarrow{\forall (id \neq f: i \rightarrow j)} & X_j \\
 \downarrow \widehat{m}_i k & \nearrow & \downarrow & & \downarrow \\
 Y_i & \xrightarrow{(Y_i \rightarrow Y_j)_{id \neq f: i \rightarrow j}} & \lim_{id \neq f: i \rightarrow j} Y_j & \xrightarrow{\forall (id \neq f: i \rightarrow j)} & Y_j
 \end{array}$$

$\widehat{M}_i k$  is positioned between  $X_i$  and  $\lim_{id \neq f: i \rightarrow j} X_j$ , with a red arrow from  $X_i$  to  $\widehat{M}_i k$  labeled  $\widehat{m}_i k$ , and a red arrow from  $\widehat{M}_i k$  to  $\lim_{id \neq f: i \rightarrow j} X_j$ . A red arrow also points from  $\widehat{M}_i k$  to  $Y_i$ . A red corner symbol  $\lrcorner$  is placed between  $\widehat{M}_i k$  and  $\lim_{id \neq f: i \rightarrow j} X_j$ .

## Definition

A category with weak equivalences is a category  $\mathbb{C}$  equipped with a **replete wide subcategory**  $\mathcal{W}_{\mathbb{C}} \hookrightarrow \mathbb{C}$  closed under **two-out-of-three** called the weak equivalences.

A homotopical diagram  $(\mathcal{I}, \mathcal{W}_{\mathcal{I}}) \rightarrow (\mathbb{C}, \mathcal{W}_{\mathbb{C}})$  is a diagram  $\mathcal{I} \rightarrow \mathbb{C}$  preserving weak equivalences.



### Example.

In the **Kapulkin–Lumsdaine** model structure on the category of models of ITT, the **path object** of each model  $\mathbb{C}$  are the **homotopical spans**  $\mathbb{C}^{0 \leftarrow 01 \rightarrow 1}$  with both legs weak equivalences.



## Definition

For finitely complete  $\mathbb{C}$ , the pushforward along  $q: B \rightarrow A \in \mathbb{C}$  is a right adjoint to the pullback along  $q$ .



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Suppose  $(\mathbb{C}, \mathcal{W}_{\mathbb{C}})$  is a category with weak equivalences equipped with a pullback-stable class of fibrations where pushforwards of fibrations along fibrations exists.



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## Theorem (Shulman 2015)

Homotopical diagrams  $\mathcal{I} \rightarrow \mathbb{C}$  closed under pushforwards of Reedy fibrations along Reedy fibrations when no maps of  $\mathcal{I}$  are weak equivalences.



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## Theorem (Kapulkin–Lumsdaine 2021)

Homotopical diagrams  $\mathcal{I} \rightarrow \mathbb{C}$  closed under pushforwards of Reedy fibrations along Reedy fibrations when all maps of  $\mathcal{I}$  are weak equivalences.



## Exponentials of spans

$$X_0 \longleftarrow X_{01} \longrightarrow X_1$$

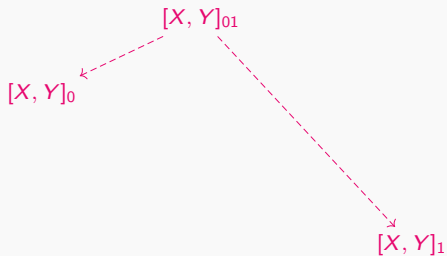
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$$[X, Y]_{01}$$



$$[X, Y]_0 \times [X, Y]_1$$



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$$\begin{array}{ccccc} X_0 & \longleftarrow & X_{01} & \longrightarrow & X_1 \\ \downarrow & & \downarrow & & \downarrow \\ Y_0 & \longleftarrow & Y_{01} & \longrightarrow & Y_1 \end{array}$$

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 \vdots & & \\
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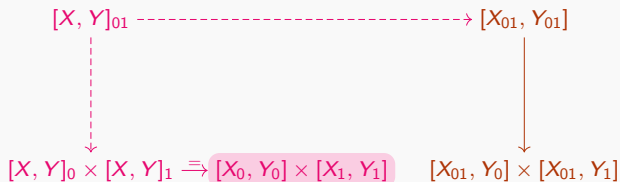
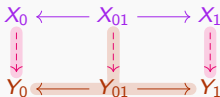
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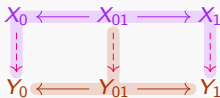


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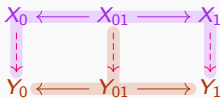


$$\begin{array}{ccc}
 [X, Y]_{01} & \xrightarrow{\quad \quad \quad} & [X_{01}, Y_{01}] \\
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 [X, Y]_0 \times [X, Y]_1 & \xrightarrow{\quad \quad \quad} & [X_{01}, Y_0] \times [X_{01}, Y_1]
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$\Rightarrow [X_0, Y_0] \times [X_1, Y_1] \rightarrow [X_{01}, Y_0] \times [X_{01}, Y_1]$



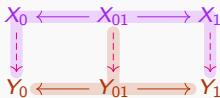
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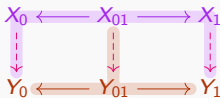
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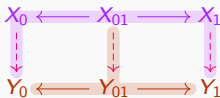
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$$\begin{array}{ccc}
 [X, Y]_{01} & \dashrightarrow & [X_{01}, Y_{01}] \\
 \downarrow & \lrcorner & \downarrow \\
 [X, Y]_0 \times [X, Y]_1 & \xrightarrow{\Rightarrow} & [X_0, Y_0] \times [X_1, Y_1] \xrightarrow{\Rightarrow} [X_{01}, Y_0] \times [X_{01}, Y_1]
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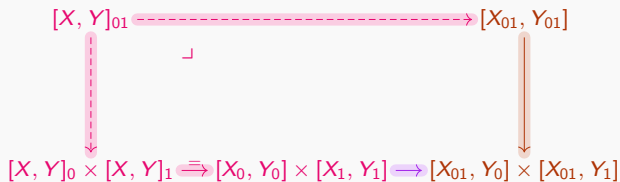
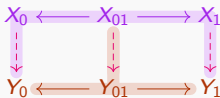
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$$\begin{array}{ccc}
 [X, Y]_{01} & \xrightarrow{\text{dashed}} & [X_{01}, Y_{01}] \\
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 [X, Y]_0 \times [X, Y]_1 & \xrightarrow{\text{dashed}} & [X_{01}, Y_0] \times [X_{01}, Y_1] \\
 \xrightarrow{\text{dashed}} & & \xrightarrow{\text{dashed}} \\
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## Exponentials of spans





## Proposition

Let  $(\mathcal{I}, \text{deg})$  be an **inverse** category and  $\mathbb{C}$  be a **cartesian closed** category. Then, for **diagrams**  $X, Y \in \mathbb{C}^{\mathcal{I}}$ , the **exponential diagram**  $[X, Y]$  is constructed by way of **induction**



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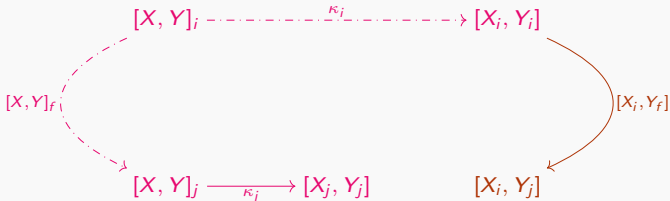


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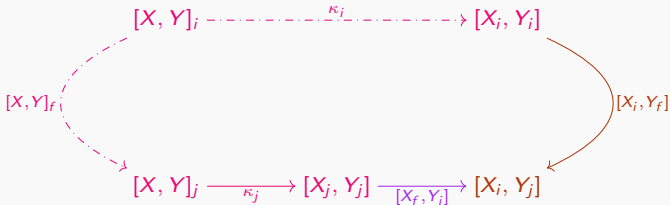


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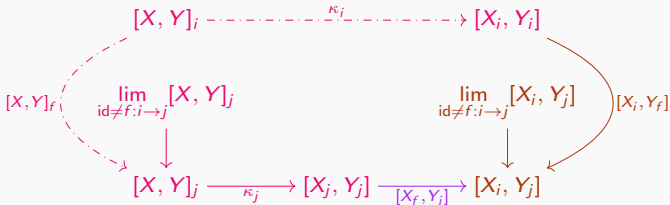


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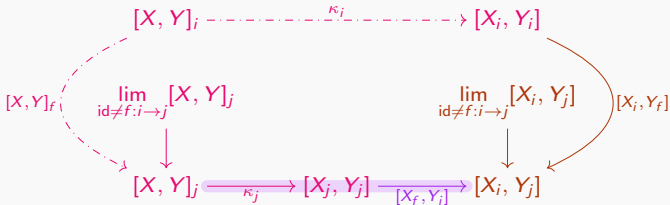


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$$\begin{array}{ccc}
 [X, Y]_i & \xrightarrow{\kappa_i} & [X_i, Y_i] \\
 \downarrow & & \downarrow \\
 \lim_{\text{id} \neq f: i \rightarrow j} [X, Y]_j & \xrightarrow{\quad} & \lim_{\text{id} \neq f: i \rightarrow j} [X_i, Y_j] \\
 \downarrow & & \downarrow \\
 [X, Y]_j & \xrightarrow{\kappa_j} [X_j, Y_j] \xrightarrow{[X_f, Y_f]} & [X_i, Y_j]
 \end{array}$$

The diagram illustrates the inductive construction of the exponential diagram  $[X, Y]$ . At the top, a dashed arrow labeled  $\kappa_i$  connects  $[X, Y]_i$  to  $[X_i, Y_i]$ . Below this, a solid arrow connects the limit  $\lim_{\text{id} \neq f: i \rightarrow j} [X, Y]_j$  to  $\lim_{\text{id} \neq f: i \rightarrow j} [X_i, Y_j]$ . A dashed arrow labeled  $[X, Y]_f$  points from  $[X, Y]_i$  to the limit  $\lim_{\text{id} \neq f: i \rightarrow j} [X, Y]_j$ . A solid arrow labeled  $[X_i, Y_f]$  points from  $[X_i, Y_i]$  to the limit  $\lim_{\text{id} \neq f: i \rightarrow j} [X_i, Y_j]$ . At the bottom, a solid arrow labeled  $\kappa_j$  connects  $[X, Y]_j$  to  $[X_j, Y_j]$ , which then connects via a solid arrow labeled  $[X_f, Y_f]$  to  $[X_i, Y_j]$ .

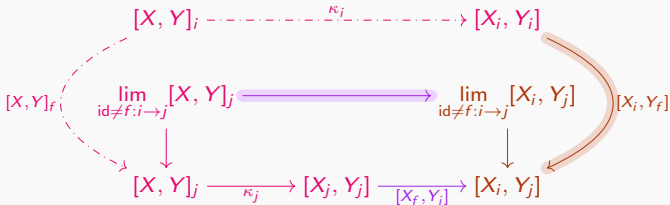


## Proposition

Let  $(\mathcal{I}, \text{deg})$  be an inverse category and  $\mathbb{C}$  be a cartesian closed category. Then, for diagrams  $X, Y \in \mathbb{C}^{\mathcal{I}}$ , the exponential diagram  $[X, Y]$  is constructed by way of induction so that at each  $i \in \mathcal{I}$ , one has a comparison map

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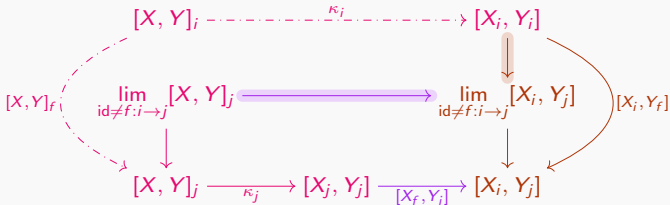


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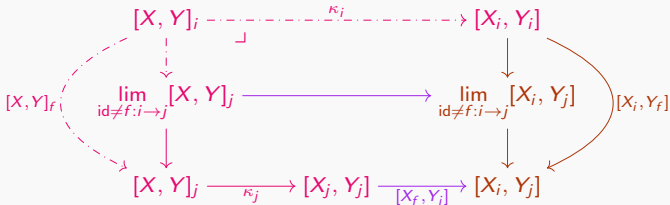


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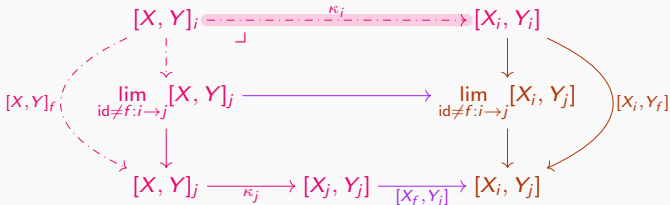


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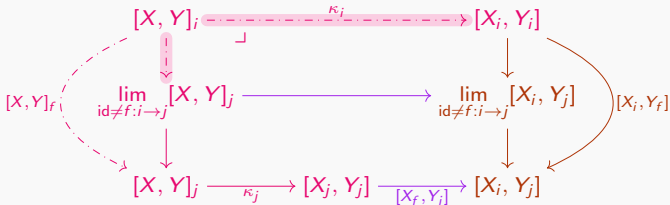


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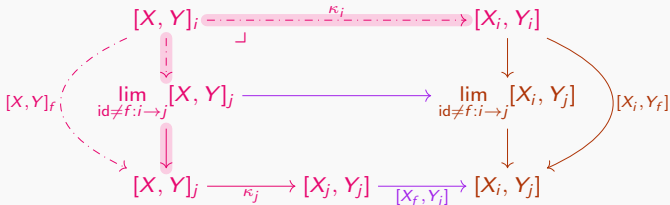


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## Assumption

- ▶ Weak equivalences of  $\mathbb{C}$  are stable under **pullback along fibrations**



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## Example.

- ▶ Models of intensional type theory such as universe categories,  $\pi$ -tribes



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## Lemma (Underslice version of Kapulkin–Lumsdaine 2021)

For all  $i \in \mathcal{I}$  if Reedy fibrant diagrams  $X, Y \in \mathbb{C}^{\mathcal{I}}$  sends all maps under  $i$  to weak equivalences then so does  $[X, Y] \in \mathbb{C}^{\mathcal{I}}$ .



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## Lemma

If  $X, Y \in \mathbb{C}^{\mathcal{I}}$  sends the initial object  $f: i \rightarrow j \in \partial(i/\mathcal{I}) = i/\mathcal{I} - \{\text{id}\}$  to a weak equivalence then so does  $[X, Y] \in \mathbb{C}^{\mathcal{I}}$ .



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## Theorem

In this case, when Reedy fibrant  $X, Y \in \mathbb{C}^{\mathcal{I}}$  are homotopical, so is  $[X, Y] \in \mathbb{C}^{\mathcal{I}}$ .



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Suppose for each  $i \in \mathcal{I}$ , if there is a weak equivalence  $\text{id} \neq w: i \rightarrow j$  then:

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Suppose for each  $i \in \mathcal{I}$ , if there is a weak equivalence  $\text{id} \neq w: i \rightarrow j$  then:

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Suppose for each  $i \in \mathcal{I}$ , if there is a weak equivalence  $\text{id} \neq w: i \rightarrow j$  then:

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## Theorem

Suppose for each  $i \in \mathcal{I}$ , if  $\text{id} \neq w: i \rightarrow j$  is a weak equivalence then

all  $i \rightarrow k$  are weak equivalences or all  $(i \rightarrow k) \neq \text{id}$  factors uniquely via  $w$

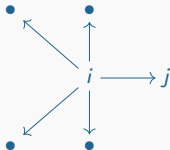
Homotopical Reedy fibrations are closed under pushforwards along homotopical Reedy fibrations.



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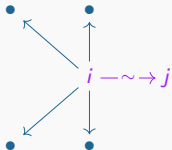




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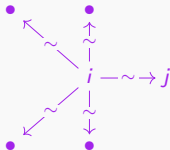




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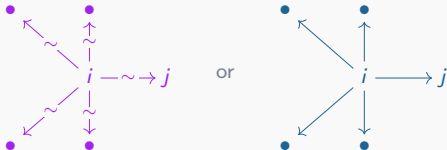




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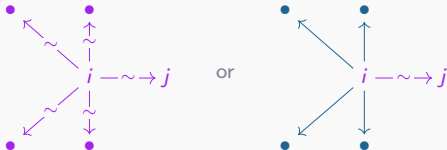
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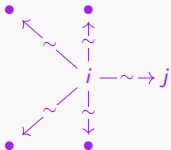
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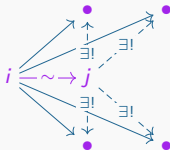
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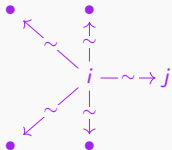
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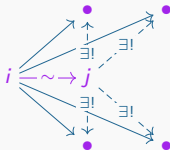
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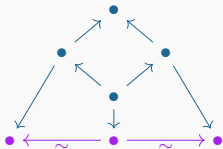


- ▶ all  $(i \rightarrow k) \neq \text{id}$  factors uniquely through  $w: i \rightarrow j$



## Example.

Valid shapes of  $\mathcal{I}$

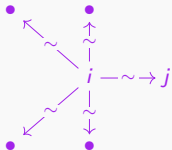




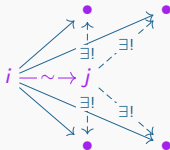
## Shape of $\mathcal{I}$

For each  $i \in \mathcal{I}$ , if  $\text{id} \neq w: i \rightarrow j$  is a weak equivalence then

- ▶ all  $i \rightarrow k$  are weak equivalences; or



or

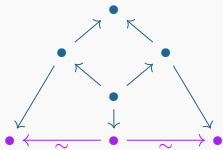


- ▶ all  $(i \rightarrow k) \neq \text{id}$  factors uniquely through  $w: i \rightarrow j$

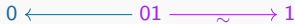


## Example.

Valid shapes of  $\mathcal{I}$



Invalid shapes of  $\mathcal{I}$

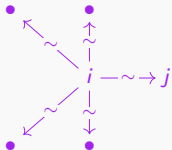




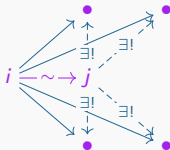
## Shape of $\mathcal{I}$

For each  $i \in \mathcal{I}$ , if  $\text{id} \neq w: i \rightarrow j$  is a weak equivalence then

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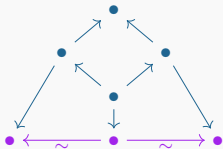


- ▶ all  $(i \rightarrow k) \neq \text{id}$  factors uniquely through  $w: i \rightarrow j$

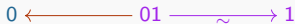


## Example.

Valid shapes of  $\mathcal{I}$



Invalid shapes of  $\mathcal{I}$



Thank you for your attention

