About the construction of simplicial and cubical sets in indexed form: the case of degeneracies

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Outline

- Reedy presheaves in (usual) fibered form vs indexed form, why?

- unary and binary parametricity as a language to uniformly talk about respectively augmented simplicial and cubical sets

- simplicial degeneracies are actually not a unary form of cubical degeneracies but of connections

- an effective uniform indexed construction of augmented simplicial and cubical sets with one degeneracy (machine-checked in Rocq)

The fibred/indexed correspondence for h-sets

For $B : \mathsf{HSet}_l$

 $\Sigma E : \mathsf{HSet}_l. (E \to B) \simeq B \to \mathsf{HSet}_l$

Application to the definition of Reedy presheaves in *indexed form*, here for cubical sets:

fibred form	vs	indexed form	
Y_0 : $HSet_l$		X_0 : $HSet_l$	(points)
Y_1 : HSet _l		$X_1 : X_0 \times X_0 \to HSet_l$	(segments)
Y_2 : $HSet_l$		$X_2 : \Pi(x_{LL}, x_{LR}) : (X_0 \times X_0) . \Pi x_{L*} : X_1(x_{LL}, x_{LR}).$ $\Pi(x_{RL}, x_{RR}) : (X_0 \times X_0) . \Pi x_{R*} : X_1(x_{RL}, x_{RR}).$	
+ coherences		$X_1(x_{LL}, x_{RL}) \times X_1(x_{LR}, x_{RR}) \to HSet_l$	(squares)
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Iterating the fibred/indexed correspondence

The domain of such X_n is a *matching object* of Y at n, say M_n , so we can build such X from such Y over a direct category by making a choice M of "matching objects". Let's write:

- c_n for the map from Y_n to M_n
- π_f for the projection to Y_p associated to the face map $f: p \to n \ (p \neq n)$
- M_f for the projection to M_p associated to the face map $f: p \to n$

Then, we have the following maps back and forth (producing back a *Reedy fibrant* Y):

fibred form

$$Y \qquad \qquad \mapsto X_n \triangleq \lambda d : M_n. \Sigma y : Y_n.(c_n(y) = d)$$

$$\begin{pmatrix} Y_n & \triangleq \Sigma d : M_n. X_n(d) \\ Y_{id} & \triangleq id \\ Y_{f:p \to n} \triangleq \lambda(d, _). (M_f(d), \pi_f(d)) \end{pmatrix} \leftrightarrow X$$

A precise study to be done, showing that we get an equivalence (not the purpose of the talk though).

Building presheaves in indexed form directly

Alternatively, we can define "matching objects" $M_n(X_0, ..., X_{n-1})$ directly on the indexed side without referring first to the fibred side. This was done, e.g., for semi-simplicial sets:

- By defining matching object as the collection of all faces, quotiented with $M_{f \circ g} = M_f \circ M_g$, as in Voevodsky 2012, Part and Luo 2015, Altenkirch, Capriotti and Kraus 2016, ...
- By relying on specific presentations of a category:
 - The d_i^n generators and $d_i d_j = d_{j-1} d_i$ coherences in H. 2013
 - By following parametricity rules in H. and Ramachandra 2025

Eventually expecting interpretations, e.g. of the universe, that more closely follow the syntax...

Iterated parametricity as a uniform approach to both augmented simplicial sets and cubical sets

It seems now established that the augmented simplicial and cubical categories only differ in the "arity" of a finite set ν :

$$\begin{array}{ll} \mathsf{Obj} & := \mathbb{N} \\ \mathsf{Hom}(p,n) & := \{l \in (\nu \sqcup \{\star\})^n \mid \mathsf{number of} \star \mathsf{in} \ l = p\} \\ g \circ f & := \begin{cases} f & \mathsf{if} \ g = \epsilon \\ a \left(g' \circ f\right) & \mathsf{if} \ g = a \ g', \mathsf{where} \ a \in \nu \\ a \left(g' \circ f'\right) & \mathsf{if} \ g = \star g', \ f = a \ f', \mathsf{where} \ a \in \nu \mathsf{ or} \ a = \star \end{cases} \\ \mathsf{id} & := \star \ldots \star \ n \ \mathsf{times} \ \mathsf{for} \ \mathsf{id} \in \mathsf{Hom}(n,n) \end{array}$$

That is, we obtain:

augmented semi-simplicial sets with $\nu = \{0\}$

semi-cubical sets with $\nu = \{L, R\}$



e.g.:
$$LR \xrightarrow{\star R} RR$$
$$L\star \mid \star \star \mid R\star$$
$$LL \xrightarrow{\star L} RL$$

Adding (one) reflexivity (in the last direction)

fibred form	vs	indexed form	
Y_0 : $HSet_l$		X_0 : $HSet_l$	(points)
Y_1 : $HSet_l$		$X_1 : X_0 \times X_0 \to HSet_l$	(segments)
$\uparrow \uparrow \uparrow \uparrow \downarrow$		$r_1 ~:~ \Pi x_0 : X_0. X_1(x_0,x_0)$	
Y_2 : $HSet_l$		X_2 : $\Pi(x_{LL}^0, x_{LR}^0) : (X_0 \times X_0) . \Pi x_{L*}^1 : X_1(x_{LL}^0, x_{LR}^0).$	
		$\Pi(x_{RL}^0, x_{RR}^0) : (X_0 \times X_0) . \ \Pi x_{R*}^1 : X_1(x_{RL}^0, x_{RR}^0).$	
+ coherences		$X_1(x_{LL}^0, x_{RL}^0) \times X_1(x_{LR}^0, x_{RR}^0) \to HSet_l$	(squares)
		r_2 : $\Pi(x_L^0, x_R^0) : (X_0 \times X_0). \ \Pi x^1 : X_1(x_L^0, x_R^0).$	
		$X_2((x^0_L, x^0_L), r_1(x^0_L), (x^0_R, x^0_R), r_1(x^0_R), (x^1, x^1))$	
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(other reflexivities can be obtained if we add also, e.g., permutations)

In terms of matching objects, reflexivities have the form:

 $r_n : \Pi d : M_n. \Pi x : X_n(d). X_{n+1}(M_r(d, x))$

for some $M_r: (\Sigma d: M_n, X_n(d)) \to M_{n+1}$ to be defined

In passing: simplicial degeneracies are *not* the unary case of cubical degeneracies but the unary case of cubical *connections*

set

reflexivities (involve one direction) *connections* (*involve two directions*)

unary $egin{array}{lll} X_{-1} : \mathsf{HSet} & & \\ X_0 : X_{-1} \to \mathsf{HSet} & & \\ X_1 : \Pi x^{-1} . \ X_0(x^{-1}) & & \\ & \to X_0(x^{-1}) \to \mathsf{HSet} & & \end{array}$

$$\begin{aligned} r_{-1} &: \Pi x^{-1} . X_0(x^{-1}) \\ r_0 &: \Pi x^{-1} . \Pi x^0 . X_1(r_{-1}(x^{-1}), x^0) \\ \text{(as in Parametric Type Theory)} \end{aligned} \qquad c_0 &: \Pi x^{-1} . \Pi x^0 . X_1(x^0, x^0) \\ \end{aligned}$$

$$\begin{array}{l} X_0:\mathsf{HSet}_l\\ X_1:X_0\times X_0\to\mathsf{HSet}_l\\ \mathsf{binary}\quad X_2:\Pi(x^0_{LL},x^0_{LR})x^1_{L\star}(x^0_{RL},x^0_{RR})x^1_{R\star}.\\ & X_1(x^0_{LL},x^0_{RL})\times X_1(x^0_{LR},x^0_{RR})\\ & \to\mathsf{HSet}_l \end{array}$$

$$egin{aligned} &r_0:\Pi x^0.X_1(x^0,x^0)\ &r_1:\Pi(x^0_L,x^0_R).\,\Pi x^1.\ &X_2((x^0_L,x^0_L),r_0(x^0_L),\ &(x^0_R,x^0_R),r_0(x^0_R),\ &(x^1,x^1)) \end{aligned}$$

(only one reflexivity per direction)

 $c_{1L} : \Pi(x_L^0, x_R^0). \ \Pi x^1. \\ X_2((x_L^0, x_L^0), r_0(x_L^0), \\ (x_L^0, x_R^0), x^1, \\ (r_0(x_L^0), x^1))$

(*n* connections per direction, completing full arity with n - 1 reflexivities)

n-ary

An effective indexed construction as a dependent stream of dependent sets

 ν -sets

$$\begin{array}{ll} \nu \mathsf{set}_l & : & \mathsf{Type}_{l+1} \\ \nu \mathsf{set}_l & \stackrel{\Delta}{=} & \nu \mathsf{set}_l^{\geq 0}(\star) \end{array}$$

$$\begin{array}{lll} \nu\mathsf{set}_l^{\geq n} & (X_{< n}:\nu\mathsf{set}_l^{< n}) & : & \mathsf{Type}_{l+1} \\ \nu\mathsf{set}_l^{\geq n} & X_{< n} & \triangleq & \Sigma X_n:\nu\mathsf{set}_l^{=n}(X_{< n}).\,\nu\mathsf{set}_l^{\geq n+1}(X_{< n},X_n) \end{array}$$

Truncated v-sets

$$\begin{array}{lll} \nu \mathsf{set}_l^{< n} & : & \mathsf{Type}_{l+1} \\ \nu \mathsf{set}_l^{< 0} & \triangleq & \mathsf{unit} \\ \nu \mathsf{set}_l^{< n'+1} & \triangleq & \Sigma X_{< n} : \nu \mathsf{set}_l^{< n'} . \, \nu \mathsf{set}_l^{= n}(X_{< n})) \end{array}$$

$$\begin{array}{lll} \nu\mathsf{set}_l^{=n} & (X_{< n}:\nu\mathsf{set}_l^{< n}) & : & \mathsf{Type}_{l+1} \\ \nu\mathsf{set}_l^{=n} & X_{< n} & \triangleq & \mathsf{fullframe}_l^n(X_{< n}) \to \mathsf{Type}_l \end{array}$$

where the "matching" fullframe $_{l}^{n}$ is defined by mutual recursive construction (see later)

The recursive process used to build frames from layers of paintings





An n-frame is made of n layers, each made of two opposite paintings of decreasing intrinsic dimension and stretched to adjust to the dimension of the frame

The recursive construction, formally

$fullframe_l^n$ $fullframe_l^n$	$\begin{array}{l} (X_{< n}:\nuset_l^{< n}) \\ X_{< n} \end{array}$: 	$Type_l$ frame $_l^{n,n}(X_{< n})$
$frame_l^{n,p,[p\leq n]}$ $frame_l^{n,0}$ $frame_l^{n,p'+1}$	$egin{aligned} &(X_{< n}: u set_l^{< n})\ &X_{< n}\ &X_{< n} \end{aligned}$: 	$egin{aligned} Type_l \ unit \ \Sigma d : frame_l^{n,p'}(X_{< n}). \ layer_l^{n,p'}(d) \end{aligned}$
$layer_l^{n,p,[p < n]}$ $layer_l^{n,p}$	$ \{X_{< n} : \nu set_l^{< n}\} \ (d : frame_l^{n, p}(X_{< n})) \\ (X_{< n-1}, X_{n-1}) \ d $:	$Type_l$ $\Pi \epsilon. painting_l^{n-1,p}(X_{n-1})(restrframe_{l,\epsilon}^{n,p}(d))$
painting $_{l}^{n,p,[p\leq n]}$ painting $_{l}^{n,p,[p=n]}$ painting $_{l}^{n,p,[p$	$ \{X_{ X_{ X_{$: ⊴∥ ⊴∥	Type _l $X_n(d)$ Σb : layer ^{n,p} _l (d). painting ^{n,p+1} _l (X _n)(d, b)

which corresponds, when $\nu = 2$, to the following organisation of the 3^n components of a *n*-cube (shown for n = 2), with frame_l associating layers on the left and painting_l associating them on the right:



The recursive construction: restrictions ("faces")

restrframe $_{m,\epsilon}^{n,q,p,[p \leq q < n]}$ restrframe $_{m,\epsilon}^{n,q,0}$	$ \{X_{< n} : \nu set_m^{< n}\} $ $ (d : frame_m^{n,p}(X_{< n})) $ $ X_{< n} \star $ $ X_{< n} (d, l) $: <u> </u> <u> </u>	frame ^{$n-1,p$} _{m} ($X_{< n-1}$) * (restriction $(X_{< n-1})$)
resumance m,ϵ	$\Lambda_{\leq n}(u,t)$	_	$(\text{resultable}_{m,\epsilon}^{-}(a), \text{resultagel}_{m,\epsilon}^{-}(l))$
	$\{X_{< n}: \nu set_m^{< n}\}$		\cdot n $1n($ n n n n n
$restrlayer_{m,\epsilon}^{n,q,p,[p < q < n]}$	$\{d: frame_m^{n,p}(X_{< n})\}$:	$layer_m^{n-1,p}(restrframe_{m,\epsilon}^{n,q,p}(d))$
n n n n n	$(t: \operatorname{layer}_{m}^{n}(a))$	Δ	$1 \qquad \xrightarrow{n p p q p (1)} (1 \qquad \cdots \qquad n - 1 q - 1 p (1))$
restriayer m,ϵ	$(X_{\leq n-1}, X_{n-1}) d l$	=	$\lambda \epsilon'$. conframe $m_{\epsilon,\epsilon'}^{m,r,r}(d)$ (restription $m_{\epsilon'}^{m,r,q}(l,q')$ ($l_{\epsilon'}$))
	$\{X_{\leq n}: \nu set_m^{\leq n}\}$		
restrocinting $n,q,p,[p \le q < n]$	$\{X_n : \nu set_m^{=n}(X_{< n})\}$		pointing $n-1, p(X)$ (restriction $n, q, p(d)$)
Test painting m, ϵ	$\{d:frame_m^{n,p}(X_{< n})\}$	•	painting _m = (X_n) (resumance _{m,e} (a))
r 1	$(c: painting_m^{n,p}(X_n)(d))$	•	
restrpainting $m, q, p, [p=q]$	$X_{< n} X_n d (l, _)$		l_{ϵ}
$restrpainting_{m,\epsilon}^{n,q,p,[p < q]}$	$X_{< n} X_n d (l, c)$	\triangleq	$(restrlayer_{m,\epsilon}^{n,q,p}(l),restrpainting_{m,\epsilon}^{n,q,p+1}(c))$

where $cohframe_{m,\epsilon,\epsilon'}$ is a coherence proof and we write $cohframe_{m,\epsilon,\epsilon'}$ for the rewriting of this proof from left to right

The recursive construction: coherences

$\begin{array}{l} cohframe_{m,\epsilon,\epsilon'}^{n,q,r,p} \\ cohframe_{m,\epsilon,\epsilon'}^{n,q,r,0} \\ cohframe_{m,\epsilon,\epsilon'}^{n,q,r,0+1} \\ cohframe_{m,\epsilon,\epsilon'}^{n,q,r,p'+1} \end{array}$	$ \begin{aligned} & \{X_{< n} : \nu set_m^{< n}\} \\ & (d : frame_m^{n, p}(X_{< n})) \\ & X_{< n} \star \\ & X_{< n} \ (d, l) \end{aligned} $: 	$\begin{split} &\operatorname{restrframe}_{m,\epsilon}^{n-1,q-1,p}(\operatorname{restrframe}_{m,\epsilon'}^{n,r,p}(d)) \\ &= \operatorname{restrframe}_{m,\epsilon'}^{n-1,r,p}(\operatorname{restrframe}_{m,\epsilon}^{n,q,p}(d)) \\ &\operatorname{refl} \star \\ &(\operatorname{cohframe}_{m,\epsilon,\epsilon'}^{n,q,r,p'}(d),\operatorname{cohlayer}_{m,\epsilon,\epsilon'}^{n,q,r,p'}(l)) \end{split}$
$cohlayer_{m,\epsilon,\epsilon'}^{^{n,q,r,p} < r < q < n]}$	$ \begin{aligned} & \{X_{< n} : \nu set_m^{< n} \{ \\ & \{d : frame_m^{n, p}(X_{< n}) \} \\ & (l : layer_m^{n, p}(X_{< n})(d)) \\ & X_{< n} \ d \ l \end{aligned} $:	$\begin{split} & restrlayer_{m,\epsilon}^{n-1,q-1,p}(restrlayer_{m,\epsilon'}^{n,r-1,p}(l)) \\ &= restrlayer_{m,\epsilon'}^{n-1,r-1,p}(restrlayer_{m,\epsilon}^{n,q,p}(l)) \\ &\lambda\epsilon''. \operatorname{cohpainting}_{m,\epsilon,\epsilon'}^{n-1,q-1,r-1,p}(l_{\epsilon''}) \end{split}$
cohpainting $m, q, r, p \ [p \le r < q < n]$ cohpainting m, ϵ, ϵ' cohpainting $m, q, r, p, [p=r]$ m, ϵ, ϵ' cohpainting $m, q, r, p, [p$	$ \begin{aligned} & \{X_{$: 	restrpainting ^{<i>n</i>-1,<i>q</i>-1,<i>p</i>} (restrpainting ^{<i>n</i>,<i>r</i>,<i>p</i>} (<i>c</i>)) = restrpainting ^{<i>n</i>-1,<i>r</i>,<i>p</i>} (restrpainting ^{<i>n</i>,<i>q</i>,<i>p</i>} (<i>c</i>)) refl (restrpainting ^{<i>n</i>-1,<i>q</i>-1,<i>p</i>} (<i>l</i> _{ϵ})) (cohlayer ^{<i>n</i>,<i>q</i>,<i>r</i>,<i>p</i>} (<i>l</i>), cohpainting ^{<i>n</i>,<i>q</i>,<i>r</i>,<i>p</i>+1} (<i>c</i>))

where we hide some coherences (such as proof-irrelevance of equality in HSet or the identification of the equality on pairs as a pair of equalities)

Adding reflexivities

We now set $\nu = 1$. For any νset_l , we define a stream of reflexivities:

$$\begin{split} \nu \mathsf{reflSet}(X_{-1}, X_0, \ldots) &\triangleq \\ \Sigma r_{-1} : \Pi d : \mathsf{frame}^{-1} . \ \Pi x : X_{-1}(d) . \ X_0(\mathsf{reflframe}^{-1}(d), x) . \\ \Sigma r_0 : \Pi d : \mathsf{frame}^0(X_{-1}) . \ \Pi x : X_0(d) . \ X_1(\mathsf{reflframe}^0(r_{-1})(d), x) . \\ \Sigma r_1 : \Pi d : \mathsf{frame}^1(X_{-1}, X_0) . \ \Pi x : X_1(d) . \ X_2(\mathsf{reflframe}^1(r_{-1}, r_0)(d), x) . \\ \ldots \end{split}$$

where

$$\mathsf{reflframe}^n(r_{-1},...,r_{n-1}):\mathsf{frame}^n(X_{-1},...,X_{n-1})\to\mathsf{frame}^{n+1,n}(X_{-1},...,X_n)$$

computes the n first layers of the border of $r_n(d)(x)$, knowing that the last layer is made of x itself, so that $(\text{reflframe}^n(r_{-1}, ..., r_{n-1})(d), x)$ is a full frame (the matching map formerly called M_r), that is of type frameⁿ⁺¹ $(X_{-1}, ..., X_n)$.

We also need two coherence conditions:

 $\begin{aligned} & \mathsf{idrestrreflframe}^n(r_{-1},...,r_{n-1}):\Pi d:\mathsf{frame}^n.\mathsf{restrframe}^{n,n}_n(\mathsf{reflframe}^n(r_{-1},...,r_{n-1})(d)) = d\\ & \mathsf{cohrestrreflframe}^n_{p < n}(r_{-1},...,r_{n-1}):\Pi d:\mathsf{frame}^{n,p}.\\ & \mathsf{restrframe}^{n,p}_p(\mathsf{reflframe}^{n,p}(r_{-1},...,r_{n-1})(d)) = \mathsf{reflframe}^{n-1,p}(r_{-1},...,r_{n-2})(\mathsf{restrframe}^{n-1,p}_p(d)) \end{aligned}$

where $reflframe^{n,p}$ generalises $reflframe^n$ to prefixes of $frame^n$:

$$\mathsf{reflframe}^{n,p}(r_{-1},...,r_{n-1}):\mathsf{frame}^{n,p}(X_{-1},...,X_{n-1})\to\mathsf{frame}^{n+1,p}(X_{-1},...,X_n)$$

Summary

- A work in progress, machine-checking a model following the iterated parametricity translation in indexed form.
- The addition of a reflexivity in the last direction is completed
- To be done: permutations, $\Pi\text{-types},$ $\Sigma\text{-types},$ universes, \dots
- Also in progress: a more compact definition relying on finer-grain dependencies between the different components of the construction.