## A note on left adjoints preserving colimits in HoTT

Perry Hart Kuen-Bang Hou (Favonia)

University of Minnesota, Twin Cities

In classical category theory, a well-known fact is that *left adjoints preserve colimits (LAPC)*. We would like to port this theorem to adjunctions between wild categories in HoTT. However, this turns out to be harder than one might hope. We isolate a sufficient condition for the proof to go through. With this condition, combined with an extension of a known technique based on homogeneous types, we show that suspension, an example of a left adjoint, preserves colimits (over graphs). Note: the sufficient condition and application to the suspension functor were already mentioned in our prior work [4, Appendix B], but here we place these results in a wider context and explain, for the first time, how we formally prove that suspension preserves colimits.

### Classical proofs of LAPC

Consider an adjunction  $L \dashv R$  between 1-categories  $\mathcal{C}$  and  $\mathcal{D}$ . Let  $\mathcal{J}$  be a small 1-category. The classical theorem states that L preserves  $\mathcal{J}$ -shaped colimits, and it has two well-known proofs. The first requires that  $\mathcal{C}$  and  $\mathcal{D}$  admit global colimit functors  $\operatorname{colim}_{\mathcal{J}} : \mathcal{C}^{\mathcal{J}} \to \mathcal{C}$  and  $\operatorname{colim}_{\mathcal{J}} : \mathcal{D}^{\mathcal{J}} \to \mathcal{D}$  [7, Chapter V.5]. The proof assumes these two colimit functors satisfy some coherence conditions that are automatically true for 1-categories (but *not* for wild ones).

Instead of requiring global colimit functors, the second proof starts with a specific colimit  $\operatorname{colim}_{\mathcal{J}}(F)$  of a diagram  $F : \mathcal{J} \to \mathcal{C}$  and shows that  $L(\operatorname{colim}_{\mathcal{J}}(F))$  is the expected colimit. Similarly to the first proof, it also assumes some coherence conditions that are true for 1-categories, and thus it is expected that we will need further assumptions when working with wild categories. Now, the proof argues that for each  $Y \in \mathsf{Ob}(\mathcal{D})$ , the following chain of isomorphisms with  $C := \operatorname{colim}_{\mathcal{J}}(F)$  equals the canonical post-composition map [8, Theorem 4.5.2]:

$$\hom_{\mathcal{D}}(L(C), Y) \cong \hom_{\mathcal{C}}(C, R(Y)) \cong \lim_{i}(\hom_{\mathcal{C}}(F_i, R(Y))) \cong \lim_{i}(\hom_{\mathcal{D}}(L(F_i), Y))$$
(iso)

This means that the induced cocone on  $L(\operatorname{colim}_{\mathcal{J}}(F))$  is indeed colimiting, i.e., L preserves colimits. Besides avoiding global colimit functors, this proof argues mostly in terms of hom-isomorphisms that can be directly supplied by a reasonable definition of adjunctions. This helps us formulate further conditions needed for wild categories. Thus, we will port the second proof to HoTT.

#### Porting to wild categories in HoTT

Let  $\mathcal{C}$  and  $\mathcal{D}$  be wild categories with functors  $L: \mathcal{C} \to \mathcal{D}$  and  $R: \mathcal{D} \to \mathcal{C}$ . Let  $\Gamma$  be a graph and  $F: \Gamma \to \mathcal{C}$  be a  $\mathcal{C}$ -valued diagram over  $\Gamma$ . Consider a cocone (C, r, K) under F where  $r_i: F_i \to C$  for vertices  $i: \Gamma_0$  and  $K_{i,j,g}: r_j \circ F_{i,j,g} = r_i$  for  $i, j: \Gamma_0$  and edges  $g: \Gamma_1(i, j)$ . As notation abuse, we write  $L(K_{i,j,g})$  as the witness for  $L(r_j) \circ L(F_{i,j,g}) = L(r_i)$  derived from the composition law  $L_\circ$  of L. Suppose this cocone (C, r, K) is colimiting. Further, suppose that  $L \dashv R$ , witnessed by a family

Suppose this cocone (C, r, K) is columiting. Further, suppose that  $L \dashv R$ , witnessed by a family of hom-equivalences  $\alpha : \prod_{X:Ob(\mathcal{D})} \prod_{A:Ob(\mathcal{C})} \hom_{\mathcal{D}}(LA, X) \xrightarrow{\simeq} \hom_{\mathcal{C}}(A, RX)$  together with proofs  $V_1$ and  $V_2$  of the naturality of  $\alpha$  in X and A, respectively. We wish to replay the second classical proof by showing the chain of isomorphisms (iso) equals the canonical post-composition map.

Let  $\zeta$  denote the composite of these isomorphisms. By function extensionality, it suffices to show  $\zeta$  and post-composition are equal on h for every  $h : \hom_{\mathcal{D}}(L(C), Y)$ . For each  $i : \Gamma_0$ , we can build an

equality  $Q_i : \alpha^{-1}(\alpha(h) \circ r_i) = h \circ L(r_i)$  from  $V_2$  and the equivalence data for  $\alpha$ . By the structure identity principle for lim, it suffices to show the following equality for all  $i, j : \Gamma_0$  and  $g : \Gamma_1(i, j)$ :

$$\mathsf{ap}_{-\circ L(F_{i,j,g})}(Q_j)^{-1} \cdot \mathsf{pr}_2(\zeta(h))(j,i,g) = \mathsf{assoc}(h, L(r_j), L(F_{i,j,g})) \cdot \mathsf{ap}_{h\circ-}(L(K_{i,j,g})) \tag{eq}$$

The problem is that this equality need not hold for wild categories.

**Example 1.** Define the wild category  $\mathcal{E}$  by  $\mathsf{Ob}(\mathcal{E}) \coloneqq *$ ,  $\hom_{\mathcal{E}}(*, *) \coloneqq S^1$ . The remaining structure on  $\mathcal{E}$  comes from the associative *H*-space structure  $\bullet$  on  $S^1$ . For all  $\ell : \hom_{\mathcal{E}}(*, *)$ , we have a nontrivial loop  $\mathsf{loop}_{\ell}$  at  $\ell$ . Let  $\Lambda : \mathcal{E} \to \mathcal{E}$  be the identity on objects and morphisms, but let  $\Lambda_{\circ}(\ell_1, \ell_2) \coloneqq \mathsf{loop}_{\ell_1 \bullet \ell_2}$ . If  $h \equiv \mathsf{id}_*$ , it is provably false that the evident hom-adjunction  $\Lambda \dashv \Lambda$  satisfies (eq).

We isolate the following general property of an adjunction as a sufficient condition for (eq).

**Definition 2.** Consider again the adjunction  $L \dashv R$ . We say that L is 2-coherent if for all  $h_1$ : hom<sub> $\mathcal{D}$ </sub>(L(X), Y),  $h_2$ : hom<sub> $\mathcal{C}$ </sub>(Z, X), and  $h_3$ : hom<sub> $\mathcal{C}$ </sub>(W, Z), the following diagram commutes:

$$\begin{array}{c} (\alpha(h_{1}) \circ h_{2}) \circ h_{3} & \xrightarrow{\text{assoc}(\alpha(h_{1}),h_{2},h_{3})} \\ \alpha(h_{1} \circ h_{2}) \circ h_{3} & \|V_{2}(h_{2} \circ h_{3},h_{1}) \\ \alpha(h_{1} \circ L(h_{2})) \circ h_{3} & \alpha(h_{1} \circ L(h_{2} \circ h_{3})) \\ V_{2}(h_{3},h_{1} \circ L(h_{2})) \| & \| \text{ap}_{\alpha}(\text{ap}_{h_{1} \circ -}(L_{\circ}(h_{2},h_{3}))) \\ \alpha((h_{1} \circ L(h_{2})) \circ L(h_{3})) & \xrightarrow{\text{assoc}(h_{1},L(h_{2}),L(h_{3})))} \alpha(h_{1} \circ (L(h_{2}) \circ L(h_{3}))) \end{array}$$

$$(2-\operatorname{coh})$$

**Theorem 3.** If L is 2-coherent, then the cocone (L(C), L(r), L(K)) under L(F) is colimiting.

*Proof.* By routine computation, (eq) is equivalent to  $(2-\cosh)$  with  $h_1$ ,  $h_2$ , and  $h_3$  instantiated with  $h, r_j$ , and  $F_{i,j,g}$ , respectively.

**Example 4.** The *n*-truncation functor  $||-||_n : A/\mathcal{U} \to A/\mathcal{U}$  is 2-coherent on all coslices of a universe. We have formalized this fact in Agda [6]. This example is critical to constructing colimits in categories of higher groups [2], which is examined in [4, Section 7.1].

#### Suspension is 2-coherent

Knowing that the suspension functor  $\Sigma : \mathcal{U}^* \to \mathcal{U}^*$  preserves colimits has a few applications. For example, this would imply that *acyclic* types [1] are closed under colimits in  $\mathcal{U}^*$ . It also would greatly simplify the construction of stable homotopy as a homology theory [3, Corollary 2.4]. Thus, we'd like to verify that  $\Sigma$  is a 2-coherent left adjoint of  $\Omega$  and thus preserves colimits. Here, the structure identity principle turns 2-coherence into a homotopy of homotopy of pointed maps.

**Definition 5.** Let  $f_1$  and  $f_2$  be pointed maps and let  $(H_1, \kappa_1), (H_2, \kappa_2) : f_1 \sim_* f_2$ . A homotopy between  $(H_1, \kappa_1)$  and  $(H_2, \kappa_2)$  consists of a homotopy  $\mu : H_1 \sim H_2$  and a path  $M_\mu : \kappa_1 =_\mu \kappa_2$  over  $\mu$ .

When building the relevant homotopy of homotopies for  $\Sigma$  in (non-Cubical) Agda, we can derive  $\mu$  from an elaborate but tractable computation. The term  $M_{\mu}$ , however, is infeasible to construct directly. Luckily, we can avoid it by adapting *Cavallo's trick* for homogeneous types [9, Homogeneous]. In fact, we use a slightly modified notion: a pointed type is *strongly homogeneous* if it's homogeneous such that the automorphism is the identity for the basepoint.

**Lemma 6.** Let  $f_1, f_2 : X_1 \to X_2$  with  $X_2$  strongly homogeneous. Let  $(H_1, \kappa_1), (H_2, \kappa_2) : f_1 \sim f_2$ . If  $H_1 \sim H_2$ , then  $(H_1, \kappa_1)$  and  $(H_2, \kappa_2)$  are homotopic.

As loop spaces are strongly homogeneous, we avoid the infeasible construction and conclude that  $\Sigma$  is 2-coherent. We have formalized this fact in Agda [5].

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