# Rezk completion of bicategories 

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(1) Rezk completion of categories
(2) Rezk completion of bicategories

- Towards computing Rezk completions
(3) Open question


## Rezk completion

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(2) Any functor $F: \mathcal{C} \rightarrow \mathcal{E}$ with $\mathcal{E}$ univalent, factors uniquely via $\mathcal{H}$ :


## Rezk completion: Definition

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such that for any univalent category $\mathcal{E}$,

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## Remark

(1) Equivalently: $\mathcal{H} \cdot(-)$ is adjoint equivalence of categories.
(2) Equivalently: $\mathcal{H} \cdot(-)$ is weak equivalence of categories.

## Rezk completion of bicategories

## Bicategories in Univalent Foundations

For every locally univalent bicategory $\mathcal{B}$, there is a univalent bicategory $\mathrm{RC}_{\text {global }}(\mathcal{B})$ and a weak equivalence $\mathcal{B} \rightarrow \mathrm{RC}_{\text {global }}(\mathcal{B})$.

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What about non-locally univalent bicategories:

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For every bicategory $\mathcal{B}$, there is a locally univalent bicategory $R C_{\text {local }}(\mathcal{B})$ and a weak equivalence $\mathcal{B} \rightarrow \mathrm{RC}_{\text {local }}(\mathcal{B})$.

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Define
(1) $\mathrm{ob}\left(\mathrm{RC}_{\text {local }}(\mathcal{B})\right):=\mathrm{ob}(\mathcal{B})$;
(2) $\mathrm{RC}_{\text {local }}(\mathcal{B})(x, y):=\mathrm{RC}(\mathcal{B}(x, y))$

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(3) What about composition?

## Rezk completion of bicategories: Composition

$$
\begin{aligned}
& \mathcal{B}(x, y) \times \mathcal{B}(y, z) \xrightarrow{\eta \times \eta} \mathrm{RC}_{\text {local }}(\mathcal{B}(x, y)) \times \mathrm{RC}_{\text {local }}(\mathcal{B}(y, z))
\end{aligned}
$$

$$
\begin{aligned}
& \mathcal{B}(x, z) \longrightarrow \mathrm{RC}_{\text {local }}(\mathcal{B}(x, z))
\end{aligned}
$$

## Rezk completion of bicategories: Left unitor



## Rezk completion of bicategories: universal property

Let $\mathcal{B}$ be a bicategory and $\eta_{\mathcal{B}}: \mathcal{B} \rightarrow \mathrm{RC}_{2}(\mathcal{B})$ be the Rezk completion of $\mathcal{B}$.

## Conjecture

For any univalent bicategory $\mathcal{D}$, the pseudo-functor

$$
\eta_{\mathcal{B}} \cdot-:\left[\mathrm{RC}_{2}(\mathcal{B}), \mathcal{D}\right] \rightarrow[\mathcal{B}, \mathcal{D}]
$$

is a bi-equivalence of bicategories. Furthermore, this characterizes the $\left(\mathrm{RC}_{2}(\mathcal{B}), \eta_{\mathcal{B}}\right)$ uniquely.
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## Rezk completion of Cat(?)

## Question

What is the Rezk completion $\mathrm{RC}_{2}$ (Cat) of Cat?

## Rezk completion of Cat: possible approach 1

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Consider the following (commuting) diagram:


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Thus $\mathrm{RC}_{2}$ (Cat) can be constructed as a displayed bicategory over Cat univ, i.e. as a bicategory of structured (univalent) categories.

## Rezk completion of Cat: possible approach 1

## Observation

Consider the following (commuting) diagram:


Thus $\mathrm{RC}_{2}$ (Cat) can be constructed as a displayed bicategory over Cat univ, i.e. as a bicategory of structured (univalent) categories.

## Question

Can we construct this displayed bicategory concretely?

## Rezk completion of Cat: possible approach 2

Another approach is:

## Question

(1) What is the local Rezk completion $\mathrm{RC}_{\text {local }}$ (Cat) of Cat?
(2) What is the Rezk completion $\mathrm{RC}_{1}([\mathcal{C}, \mathcal{D}])$ of a functor category $[\mathcal{C}, \mathcal{D}]$ ?

## Understanding the Rezk completion

## Proposition

For every $\mathcal{C}$ and $\mathcal{D}$, TFAE;
(1) $\mathrm{RC}([\mathcal{C}, \mathcal{D}])=[\mathcal{C}, \mathrm{RC}(\mathcal{D})]$;
(2) $\mathcal{C}$ and $\mathcal{D}$ are equivalent to a univalent category.

## Understanding the Rezk completion

## Proposition

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## Lemma

Let $\mathcal{B}$ be a full sub-bicategory of Cat. TFAE:
(1) $\prod_{\mathcal{C}, \mathcal{D}: \mathcal{B}} \mathrm{RC}_{1}([\mathcal{C}, \mathcal{D}])=\left[\mathcal{C}, \mathrm{RC}_{1}(\mathcal{D})\right]$;
(2) $\mathrm{RC}_{2}(\mathcal{B})=\mathcal{B}_{\text {univ }}$ (where $\mathcal{B}_{\text {univ }}$ is the intersection of $\mathcal{B}$ and Cat ${ }_{\text {univ }}$ ) ;
(3) $\mathcal{B} \simeq \mathcal{B}_{\text {univ }}$.

## Rezk completion of Cat

## Corollary

$\mathrm{RC}($ Cat $) \neq$ Cat $_{\text {univ }}$.

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## Intuition 1

(1) Cat ${ }_{\text {univ }} \sim$ categories up to weak equivalence (unique representing object) ;
(2) $\mathrm{RC}($ Cat $) \sim$ categories up to isomorphism of isomorphic functors.

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## Corollary

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## Intuition 2

(1) RC: Cat $\rightarrow \mathbf{C a t}_{\text {univ }}$ is objectwise free;
(2) $\eta_{\mathcal{C}}:$ Cat $\rightarrow \mathrm{RC}(\mathbf{C a t})$ is (externally) free

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## Question

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## Question

Is there a completion GC on bicategories such that
(1) $\mathrm{GC}($ Cat $)=$ Cat $_{\text {univ }}$;
(2) $\mathrm{GC}($ MonCat $)=$ MonCat $_{\text {univ }}$;
(3) GC(DagCat $)=$ DagCat $_{\dagger-\text { univ }}$ ?

## What does such a completion mean?

## Question

Is it possible (universally) characterize $\mathrm{RC}: \mathbf{C a t} \rightarrow \mathbf{C a t}_{\text {univ }}$ ?

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Is it possible (universally) characterize $\mathrm{RC}: \mathbf{C a t} \rightarrow \mathbf{C a t}_{\text {univ }}$ ?
This question is closely related to:
Question
Is it possible to characterize the correct notion of univalence of (structured) categories?

Thank you! Any questions?

