

\mathbb{N} from \mathbb{Z}

Christian Sattler & David Wärn

Terminology

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- ▶ a self-map $S : \mathbb{N} \rightarrow \mathbb{N}$.

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\mathbb{Z} is any type freely generated by:

- ▶ an element $0 : \mathbb{Z}$,
- ▶ a self-*equivalence* $S : \mathbb{Z} \rightarrow \mathbb{Z}$.

\mathbb{Z} -induction

Given

- ▶ $P : \mathbb{Z} \rightarrow \mathbb{U}$,
- ▶ $0_P : P(0)$,
- ▶ $S_P : (x : \mathbb{Z}) \rightarrow P(x) \simeq P(S(x))$,

obtain

- ▶ $p : (x : \mathbb{Z}) \rightarrow P(x)$,
- ▶ $p(0) = 0_P$,
- ▶ $p(S(x)) = S_P(p(x))$.

Our result

Our setting is type theory with:

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- ▶ **2** with large elimination (or descent)

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Theorem (Sattler and W.)

Given \mathbb{Z} , can construct \mathbb{N} .

We have two proofs. We present one of them in this talk.

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Thus $\Omega S^1 = \mathbb{Z}$.

Previous work

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- ▶ together with $*$: $\mathbf{1}$ this gives $\mathbb{Z} \rightarrow \mathbb{N} + \mathbf{1} + \mathbb{N}$.

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This characterizes *negative*, *zero*, and *positive* integers.

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Lemma

$$\mathbb{Z} + \mathbb{Z} \simeq \mathbb{Z}.$$

Proof.

Via doubling and halving (direct integer induction). □

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This induces:

- ▶ $\mathbb{Z} \rightarrow \mathbb{Z} + \mathbb{Z} + \mathbb{Z}$,
- ▶ hence $\text{sign} : \mathbb{Z} \rightarrow \mathbf{1} + \mathbf{1} + \mathbf{1}$,
- ▶ hence a decomposition $\mathbb{Z} \simeq \mathbb{Z}^- + \mathbb{Z}^0 + \mathbb{Z}^+$
with $S(x) \in \mathbb{Z}^+$ iff $x \in \mathbb{Z}^0 + \mathbb{Z}^+$.

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Aim: define \mathbb{N} as Σ -type over $M \equiv_{\text{def}} \mathbb{Z}^0 + \mathbb{Z}^+$.

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Both

- ▶ definition of \mathbb{N}
- ▶ derivation of \mathbb{N} -induction

use the idea of *partially defined inductive functions*.

Ordering

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Take $x \leq y$ to mean $x - y \in \mathbb{Z}^- + \mathbb{Z}^0$.

Have:

- ▶ if $x \leq y$ then $x \leq S(y)$
- ▶ if $S(x) \leq y$ then $x \leq y$
- ▶ $S(x) \leq S(y)$ iff $x \leq y$
- ▶ $x \leq 0$ iff $x \in \mathbb{Z}^0$
- ▶ $x \leq x$

Inductive functions

Let $A : M \rightarrow U$ with:

- ▶ $0_A : (x : \mathbb{Z}^0) \rightarrow A(x)$
- ▶ $S_A : (x : M) \rightarrow A(x) \rightarrow A(S(x))$

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Define $B : M \rightarrow U$ as

$$B(u) \equiv_{\text{def}} (x : M) \rightarrow x \leq u \rightarrow A(x).$$

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Have canonical maps:

- ▶ $\text{res}_u : B(S(u)) \rightarrow B(u)$
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Say $f : B(u)$ is *inductive* if $\text{res}_u(\text{ext}_u(f)) = f$.

Write $I(u)$ for type of inductive functions.

Rolling rule

For $t : X \rightarrow X$, let $\text{fix}(f) \equiv_{\text{def}} (x : X) \times (f(x) = x)$.

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Lemma

$\text{fix}(f \circ g) \simeq \text{fix}(g \circ f)$ for $f : X \rightarrow Y$, $g : Y \rightarrow X$.

Proof.

Both types are equivalent to

$(x : X) \times (y : Y) \times (f(x) = y) \times (g(y) = x)$.



Inductive functions (cont.)

$$\begin{aligned} I(0) &\simeq \text{fix}(\text{res}_0 \circ \text{ext}_0) \\ &\simeq \text{fix}(- \mapsto 0_A) \\ &\simeq \mathbf{1} \end{aligned}$$

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$$\begin{aligned} I(S(u)) &\simeq \text{fix}(\text{res}_{S(u)} \circ \text{ext}_{S(u)}) \\ &\simeq \text{fix}(\text{ext}_u \circ \text{res}_u) \\ &\simeq \text{fix}(\text{res}_u \circ \text{ext}_u) \\ &\simeq I(u) \end{aligned}$$

Defining \mathbb{N}

Instantiate A as follows:

$$A(-) \equiv_{\text{def}} M$$

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$$S_A(x) \equiv_{\text{def}} S(x)$$

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Then define naturals:

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Have:

- ▶ unique element $(0, 0_N)$ of $(m : \mathbb{Z}^0) \times N(m)$
- ▶ $S_N : (m : M) \rightarrow N(m) \simeq N(S(m))$

Deriving \mathbb{N} -induction

Given:

- ▶ $P : (m : M) \rightarrow N(m) \rightarrow U$
- ▶ $0_P : P(0, 0_N)$
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Prove $(m : M) \rightarrow I(m)$ by \mathbb{Z} -induction.

Deduce $(m : M) (n : N(m)) \rightarrow P(m, n)$ compatible with 0_P and S_P .